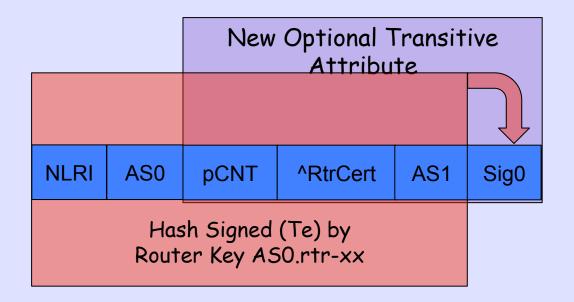
Estimating CPU Cost of BGPsec on a Router

Cisco NAG / San Jose

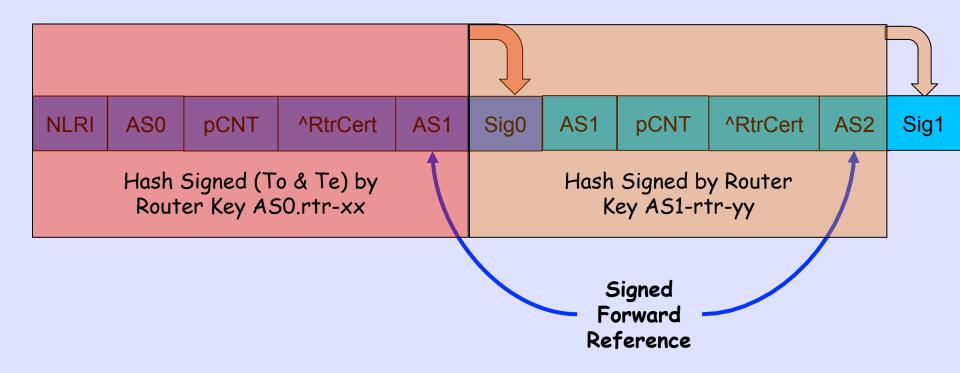
2011.10.26

Kotikalapudi Sriram <kotikalapudi.sriram@nist.gov>
Randy Bush <randy@psg.com>

BGPsec from ASO to AS1



BGPsec AS1 to AS2

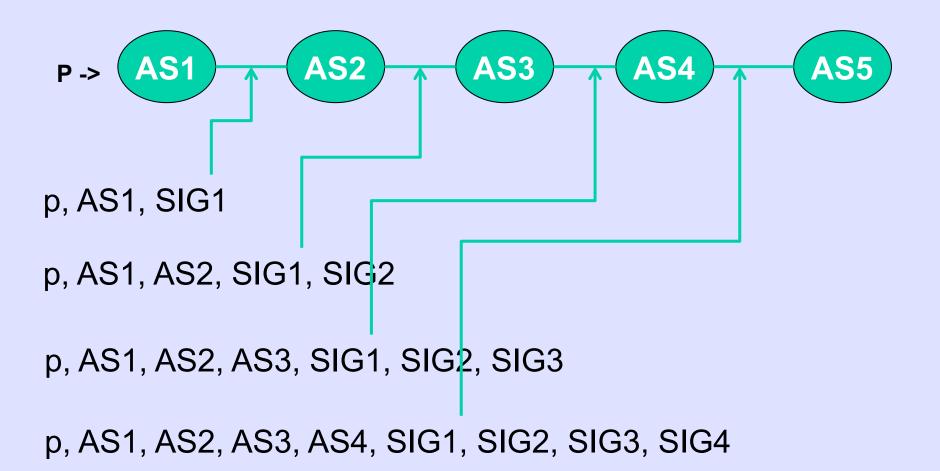


R1 signing over R0's signature is same as signing over entire R0 announcement

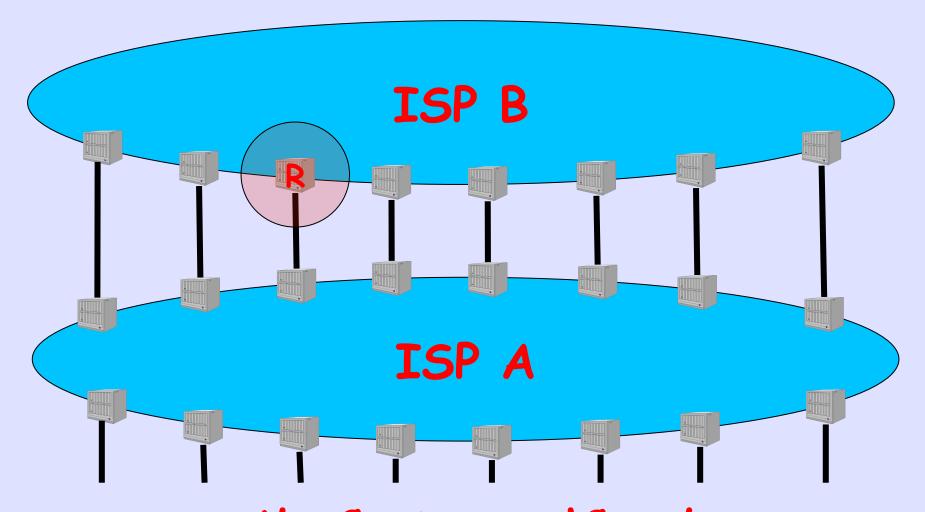
BGPsec Islands

- RPKI-Based Origin Validation can be deployed by randomly scattered ISPs
- · Each gets the benefit of origin validation
- · BGPsec depends on your neighbor signing
- It will deploy as islands which eventually interconnect

We Draw Pictures Like This



But Reality is This



A's Customer 'Cone'

Number of Paths

- One ISP router, R, has many paths for prefix P
- All but one are from iBGP peers
- BGPsec spec says R does not validate paths received from iBGP peers
- I.e. R has to validate only one path for each P from peer A

Some Largish ISPs Cones

Very Large Global

- 1 1353 --- **ISP's Own Pf**x
- 2 21586 --- **BGP Cust Pfx**
- 3 6820 --- Cust's Cust Pfx
- 4 1627 --- ...
- 5 942
- 6 45
- 7 14
- 8 6

Very Large Global

- 1 620
- 2 16028
- 3 9434
- 4 2922
- 5 435
- 6 46
- 7 15
- 8 27
- 9 1

Large Global

- 1 443
- 2 8197
- 3 8052
- 4 2715
- 5 387
- 6 37
- 7 48
- 8 157
- 9 2

Large Global

- 1 501
- 2 3686
- 3 3603
- 4 816
- 5 45
- 6 9
- 8 1

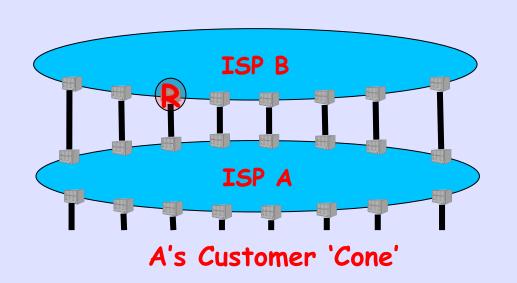
Asian Regional

- 1 152
- 2 791
- 3 120
- 4 35
- 5 3 # pfxs path length

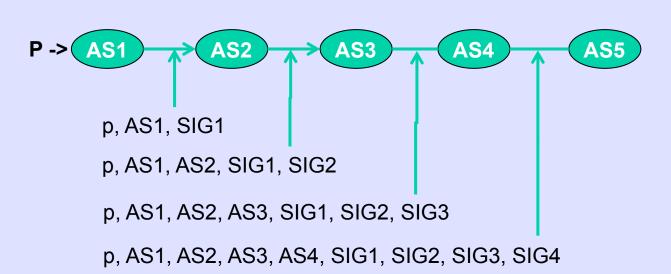
Yes, there are rather long tails

Yes, we removed prepending

Incremental Deployment



If A and B Deploy BGPsec, What is the Load on a Router?



Now this Picture Makes Sense!

Cost to Sign/Validate Using One Core

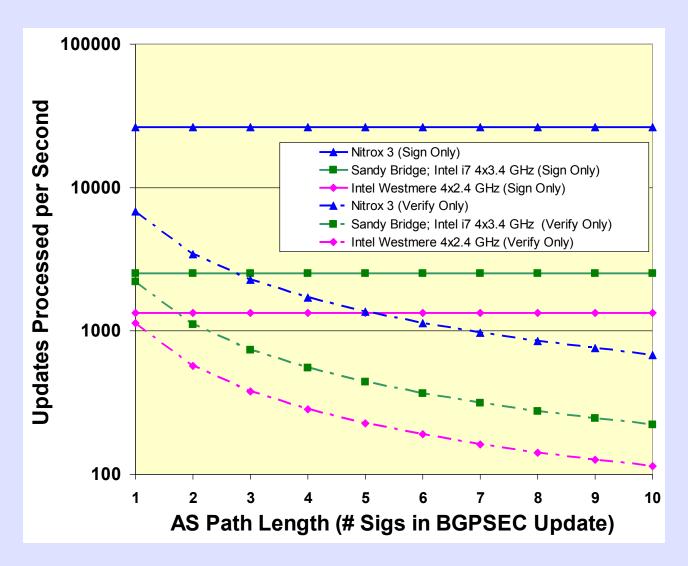
| | Operations per second | | | | |
|-------------------|-----------------------|---------------------|--------------|-------------------|------------------|
| | | | amd64, Sandy | | |
| | | | Bridge; 2011 | | |
| | | amd64; Westmere | Intel i7- | NITROX PX PCI- | NITROX III PCI- |
| | Intel Core 2 Duo, | (206c2); 2010 Intel | 2600K; 4 x | Express CN1620 - | Express CNN3570- |
| | 64-bit, 3 GHz, | Xeon E5620; 4 x | 3400MHz; | PCle Look-aside | PCIe Look-aside |
| | 8GB, Linux 5.7 | 2400MHz | threads | P rocessor | Processor |
| ECDSA-P256 Verify | 890 | 1139 | 2215 | 854 | 6832 |
| ECDSA-P256 Sign | 1100 | 1335 | 2530 | 3293 | 26344 |

 Source: eBACS: ECRYPT Benchmarking of Cryptographic Systems

http://bench.cr.yp.to/results-sign.htm

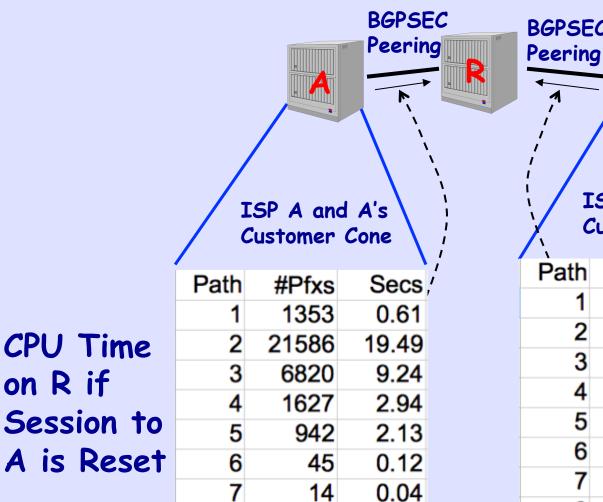
And: Cavium, Inc. (private communication)

Updates Per Second



Validation Cost Model

BGPSEC



0.02

34.59

Total Seconds

ISP C and C's Customer Cone Path #Pfxs Secs 620 0.28 16028 9434 12.78 5.28 2922 435 0.98 46 15 0.05 27 0.10 9 0.00 Total Seconds 34.06

14.47 CPU Time on R if Session to 0.12 C is Reset

2011.10.26 BGPsec Crypto Cost

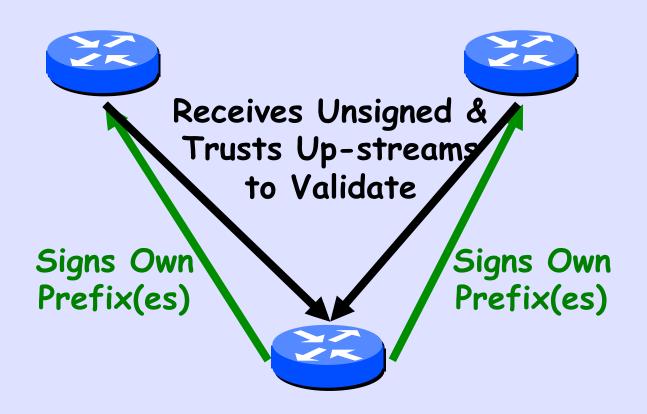
CPU Time

on R if

Signing Cost

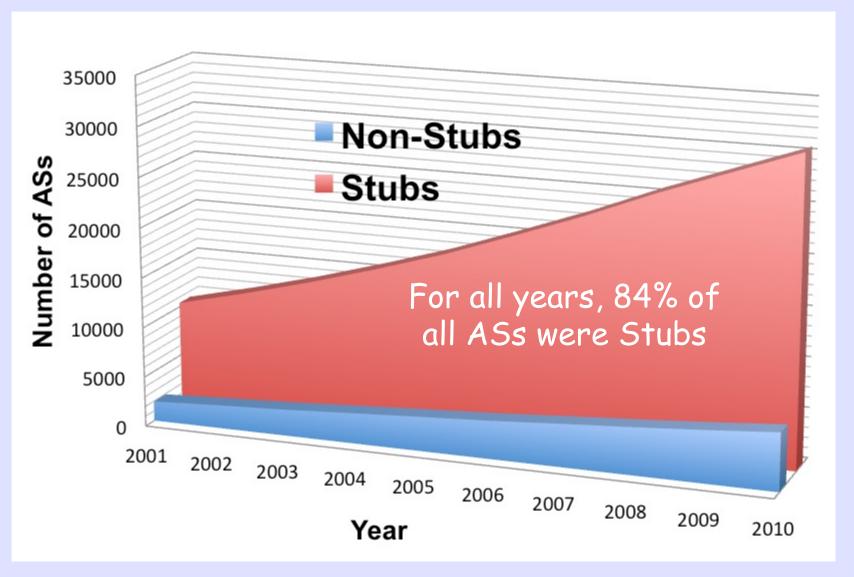
- You only sign once, irrespective of path length
- You only sign toward BGPsec speakers
- Though the cost of stripping
 BGPsec toward non-speakers may
 be on the order of signing

Need not Sign To Stubs



Only Needs to Have Own Private Key, No Other Crypto or RPKI Data No Hardware Upgrade!!

Stub ASs vs Transit



BGP Peers per Router

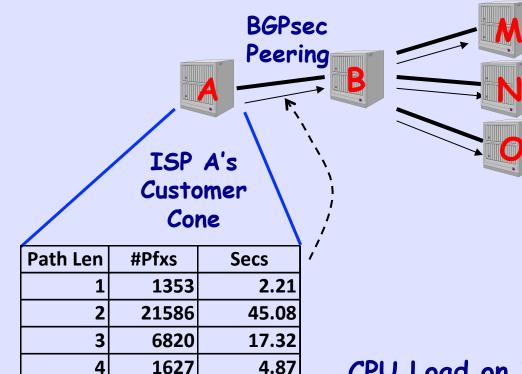
| ISP | BGP Peers | BGP Custs |
|-----|-----------|-----------|
| W | 29 | 95 |
| X | 3-4 | 20 |
| У | 6 | 12 |
| Z | 8 | 16 |

These numbers are from real ISPs, but large ones

Signing Bottom Line

- Except for W, it comes to 2-3 BGPsec customers per aggregation router
- Say 400k routes at 2530 sigs/sec
- (3*400000)/2530 = 475 seconds
- But this presumes the entire Internet is signed, which is a looooooooong time from now
- · But W will eventually have a problem!

CPU for Validation and Signing



3.24

0.18

0.06

0.03

73

942

45

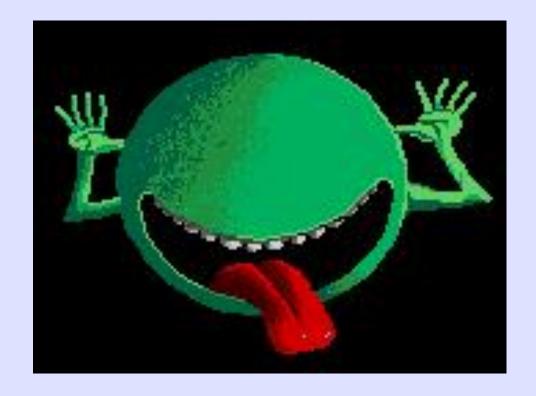
14

Total (seconds)

6

- B peers with four BGPsec peers
- B's other peers are not BGPsec aware

CPU Load on B, including Validation & Signing, if Session to A is Reset.



So Don't Panic, Engineer Prudently